

Computable Aspects of Inner Functions

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Outline

- 1 Results on analytic functions
- 2 Results on bounded analytic functions
 - Background from analysis
 - Computability results
- 3 References



The aim of this talk We will discuss some results in computable complex analysis.

Let f, g, h range over analytic functions only.



Naming systems

- $\theta_{<}$ for set of open subsets of \mathbb{C} .
- ψ for set of closed subsets of \mathbb{C} .
- $[\rho^2]^\omega$ for set of countably infinite sequences of complex numbers.
- $[\rho^2 \rightarrow \rho^2]_{\mathbb{D}}$ for the set of continuous $f : \mathbb{D} \rightarrow \mathbb{C}$
- δ_{open} for set of continuous $f : \subseteq \mathbb{C} \rightarrow \mathbb{C}$ with open domain.
 This representation is defined by:

$$\delta_{open}(\langle p, q \rangle) = f \Leftrightarrow \eta_p^{\omega\omega}(\rho^2, \rho^2)\text{-realizes } f \wedge \theta_{<}(q) = \text{dom}(f).$$



Nice properties of analytic functions

- Infinitely differentiable
- Cauchy integral formula for derivatives:
$$f'(z_0) = \frac{1}{2\pi i} \int_{\gamma} \frac{f(z)}{z-z_0} dz$$
- Cauchy's theorem: $\frac{1}{2\pi i} \int_{\gamma} \frac{f'}{f} = Z - P.$



These nice analytic properties can be used to obtain nice computability properties. e.g.

Theorem

(Hertling, 1999) $f \mapsto f'$ is $(\delta_{open}, \delta_{open})$ -computable.

Proof uses Cauchy Integral Formula.



Results on zero-sets Let f range over analytic functions that are not identically zero.

Theorem

Computable Zero-Finding (Matheson, McNicholl 2006) *The map*

$$f \mapsto f^{-1}[\{0\}]$$

is (δ_{open}, ψ) -computable. Hence, the zero set of a δ_{open} -computable analytic function is ψ -computable.

Proof uses Cauchy's Theorem.



Zero sequences

Definition

If u is analytic, then $\{a_n\}_{n=0}^{\infty}$ is a *primary zero sequence* of u if its terms are precisely the zeros of u and the number of times each zero of u is repeated is its multiplicity.

Definition

If u is analytic, then $\{a_n\}_{n=0}^{\infty}$ is a *truncated zero sequence* of u if its terms are precisely the zeros of u *besides zero*, and the number of times each zero of u is repeated is its multiplicity.



Let g range over analytic functions with infinitely many zeros, but not identically zero.

Theorem

Computable primary zero sequence (Matheson, McNicholl 2006) *There is a $(\delta_{open}, [\rho^2]^\omega)$ -computable function Ψ such that $\Psi(g)$ is a primary zero sequence of g for all g .*

Corollary

There is a $(\delta_{open}, [\rho^2]^\omega)$ -computable function Ψ such that $\Psi(g)$ is a truncated zero sequence of g for all g .



Throughout the rest of this talk, we will focus on bounded analytic functions with simply connected domain.

In addition to their analytic properties, these functions have nice *algebraic* properties which can be used to get computability results. In addition, these functions form the mathematical foundation for control theory in engineering.



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The class $H^\infty(\mathbb{D})$

- $H^\infty(\mathbb{D})$ is the set of all bounded analytic functions $f : \mathbb{D} \rightarrow \mathbb{C}$.
- For $f \in H^\infty(\mathbb{D})$, let

$$\|f\|_\infty = \sup\{|f(z)| : z \in \mathbb{D}\}.$$

- $H^\infty(\mathbb{D})$ is a Banach space under $\|\cdot\|_\infty$.



Definition

$u \in H^\infty(\mathbb{D})$ is *inner* if $\lim_{z \rightarrow z_0} |u(z)| = 1$ for almost all $z_0 \in \partial\mathbb{D}$.

Inner functions “generate” $H^\infty(\mathbb{D})$. Two important classes of inner functions:

- Singular functions
- Blaschke products



Definition

A function $s \in H^\infty(\mathbb{D})$ is *singular* if there is a finite positive Borel measure on $\partial\mathbb{D}$, μ , that is singular with respect to Lebesgue measure and such that

$$s(z) = \exp \left\{ - \int_{-\pi}^{\pi} \frac{e^{it} + z}{e^{it} - z} d\mu(t) \right\}$$

Theorem

If s is singular, then s is inner, s has no zeros, and $s(0)$ is a positive real number.



Definition

Let $A = \{a_n\}_{n=0}^{\infty}$ be a sequence of points in $\mathbb{D} - \{0\}$. The product

$$B_{A,k}(z) =_{df} z^k \prod_{n=0}^{\infty} \frac{|a_n|}{a_n} \frac{a_n - z}{1 - \overline{a_n}z}$$

is called a *Blaschke product*. We abbreviate $B_{A,0}$ with B_A .



Definition

Let $A = \{a_n\}_{n=0}^{\infty}$ be a sequence of points in $\mathbb{D} - \{0\}$. The series

$$\Sigma_A =_{df} \sum_{n=0}^{\infty} (1 - |a_n|)$$

is called the *Blaschke sum* of A . The inequality $\Sigma_A < \infty$ is called the *Blaschke condition*.



Theorem

Let $A = \{a_n\}_{n=0}^{\infty}$ be a sequence of points in $\mathbb{D} - \{0\}$.

- 1 If A satisfies the Blaschke condition, then $B_{A,k}$ is an inner function.
- 2 If A satisfies the Blaschke condition, then the terms of A are precisely the zeros of B_A . Furthermore, the number of times a zero of B_A appears in A is its multiplicity.
- 3 If A does not satisfy the Blaschke condition, then $B_A \equiv 0$.



Importance of Blaschke products

- Factorization
- Estimation of inner functions (Frostman's Theorem)



Theorem

(Factorization of Inner Functions) *If u is an inner function, then there exist unique λ_u, b_u, s_u such that $u = \lambda_u b_u s_u$, $\lambda_u \in \partial\mathbb{D}$, b_u is a (possibly finite) Blaschke product, and s_u is a singular function.*



Capacity For each closed $K \subseteq \mathbb{D}$ and each positive measure σ on K , let $U_\sigma : \mathbb{D} \rightarrow \mathbb{D}$ be defined by the equation

$$U_\sigma(z) = \int_K \log \frac{1}{|z - \zeta|} d\sigma(\zeta).$$



Definition

Let $F \subseteq \mathbb{D}$ be closed. We say that F has *zero capacity* if for every positive measure on F , σ , with $\sigma \neq 0$, U_σ is not bounded on any neighborhood of F . Otherwise, we say that F has *positive capacity*. If U is an arbitrary subset of \mathbb{D} , then we say that U has positive capacity just in case it has a closed subset with positive capacity; otherwise, we say that it has zero capacity.



Theorem

Every zero-capacity set has measure zero.

However, the Cantor middle-third set has *positive* capacity.



For $a, z \in \mathbb{D}$ with $|a| < 1$, let

$$M_a(z) = \frac{z - a}{1 - \bar{a}z}.$$

Theorem

(Frostman's Theorem) *Let u be a non-constant inner function. Then, $M_a \circ u$ is a unit multiple of a Blaschke product for all $a \in \mathbb{D}$ except in a set of capacity zero.*

The set of values of a for which $M_a \circ u$ is not a unit multiple of a Blaschke product is called the *exception set* of u .



As $a \rightarrow 0$, $\| M_a \circ u - u \|_\infty \rightarrow 0$.

Corollary

If u is a non-constant inner function, and if $\epsilon > 0$, then there is a unit multiple of a Blaschke product B such that $\| u - B \|_\infty < \epsilon$.



Theorem

A truncated zero sequence of an inner function is a Blaschke sequence.

Definition

If u is an inner function, then let Σ_u denote the Blaschke sum of its truncated zero sequence.

If u is an analytic function, then let k_u denote the order of u 's zero at 0 if there is one; if $u(0) \neq 0$, then let $k_u = 0$.



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Computable primary zero sequence (Matheson, McNicholl 2006) *There is a $(\delta_{open}, [\rho^2]^\omega)$ -computable function Ψ such that $\Psi(g)$ is a primary zero sequence of g for all g .*

Corollary

There is a $(\delta_{open}, [\rho^2]^\omega)$ -computable function Ψ such that $\Psi(g)$ is a truncated zero sequence of g for all g .



Theorem

Computable Frostman Theorem (McNicholl, 2007) *There is a $([\rho^2 \rightarrow \rho^2]_{\mathbb{D}}, \nu_{\mathbb{N}}, \nu_{\mathbb{N}})$ -computable function Ψ such that if u is inner (and non-constant) and $n \in \mathbb{N}$, then $M_{\Psi(u,n)} \circ u$ is a unit multiple of a Blaschke product and $\|u - M_{\Psi(u,n)} \circ u\|_{\infty} < 2^{-n}$.*

This theorem shows we can *effectively* estimate inner functions by unit multiples of Blaschke products.



Computability of k_u Throughout the remaining slides, u ranges over inner functions with infinitely many zeros.

Theorem

(McNicholl, 2007) $u \mapsto k_u$ is not $([\rho^2 \rightarrow \rho^2]_{\mathbb{D}}, \nu_{\mathbb{N}})$ -computable.

Theorem

(McNicholl, 2007) $(u, \Sigma_u) \mapsto k_u$ is $([\rho^2 \rightarrow \rho^2]_{\mathbb{D}}, \rho, \nu_{\mathbb{N}})$ -computable.



Theorem

(Matheson, McNicholl, 2006) *There is a $[\rho^2]^\omega$ -computable sequence $A = \{a_n\}_{n=0}^\infty$ such that B_A is not (ρ^2, ρ^2) -computable.*

In other words, merely knowing the Blaschke sequence is not enough to compute the Blaschke product.



Theorem

(Matheson, McNicholl 2006) *The map*

$$(A, \Sigma_A) \mapsto B_A$$

is $([\rho^2]^\omega, \rho, [\rho^2 \rightarrow \rho^2]_{\mathbb{D}})$ -computable.

In other words, if you know a Blaschke sequence and its Blaschke sum, then you can compute the Blaschke product.



Theorem

(McNicholl, 2007) *The map $(A, B_A) \mapsto \sum_A$ is $([\rho^2]^\omega, [\rho^2 \rightarrow \rho^2]_{\mathbb{D}}, \rho^2)$ -computable. In fact, $(A, B_A(0)) \mapsto \sum_A$ is $([\rho^2]^\omega, \rho, \rho^2)$ -computable.*

In other words, once you know a Blaschke sequence, in order to compute the Blaschke product you have to know the Blaschke sum (or an equivalent piece of information).



Corollary

(McNicholl 2007) *Suppose A is $[\rho^2]^\omega$ -computable. If B_A maps ρ^2 -computable complex numbers to ρ^2 -computable complex numbers, then B_A is (ρ^2, ρ^2) -computable.*

This is not the case for power series!

Theorem

(Caldwell, Pour-El 1975) *There is a $[\rho]^\omega$ -computable sequence $\{a_n\}_{n=0}^\infty$ such that $z \mapsto \sum_{n=0}^\infty a_n z^n$ is (ρ^2, ρ^2) -computable on closed disks but not on the entire plane.*



Theorem

(McNicholl, 2007) *The map $u \mapsto (\lambda_u, b_u, s_u)$ is not $([\rho^2 \rightarrow \rho^2]_{\mathbb{D}}, \rho^2, [\rho^2 \rightarrow \rho^2]_{\mathbb{D}}, [\rho^2 \rightarrow \rho^2]_{\mathbb{D}})$ -continuous.*

In other words, merely knowing an inner function is not enough to compute its factorization.



Theorem

(McNicholl, 2007) *The map $(u, \sum_u) \mapsto (\lambda_u, b_u, s_u)$ is $([\rho^2 \rightarrow \rho^2]_{\mathbb{D}}, \rho, \rho^2, [\rho^2 \rightarrow \rho^2]_{\mathbb{D}}, [\rho^2 \rightarrow \rho^2]_{\mathbb{D}})$ -computable.*

In other words, for inner functions with infinitely many zeros, knowing the Blaschke sum of u is *sufficient* for computing the factorization of u .







Theorem

(McNicholl, 2007) *The map $(u, b_u) \mapsto \sum_u$ is $([\rho^2 \rightarrow \rho^2]_{\mathbb{D}}, [\rho^2 \rightarrow \rho^2]_{\mathbb{D}}, \rho)$ -computable.*

In other words, to compute the factorization of u , knowledge of the Blaschke sum of u (or equivalent information) is *necessary*.



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